Gorillaz
“Clint Eastwood”
(Gorillaz)

I’m happy, I’m feeling glad
I got a new toolbox, in my bag
My code’s useless, but not for long
I know what my tasks de-pend on
I’m happy, I’m feeling glad
I used the 418 skills, in my bag
My code is per-formant
and structured for parallelization

parallelization…
parallelization…
parallelization…
Quiz

This is an ISPC function.

It contains a loop nest.

Which iterations of the loop(s) are parallelized by ISPC? Which are not?

Answer: none
Creating a parallel program

- Thought process:
  1. Identify work that can be performed in parallel
  2. Partition work (and also data associated with the work)
  3. Manage data access, communication, and synchronization

- Recall one of our main goals is speedup *
  For a fixed computation:

\[
\text{Speedup( P processors )} = \frac{\text{Time (1 processor)}}{\text{Time (P processors)}}
\]

* Other goals include high efficiency (cost, area, power, etc.) or working on bigger problems than can fit on one machine
Creating a parallel program

Problem to solve → Decomposition

Subproblems (a.k.a. “tasks”, “work to do”)

Parallel Threads ** (“workers”)

Parallel program (communicating threads)

Execution on parallel machine

Decomposition → Assignment

Parallel Threads **

Orchestration

Parallel program (communicating threads)

Mapping

These responsibilities may be assumed by the programmer, by the system (compiler, runtime, hardware), or by both!

Adopted from: Culler, Singh, and Gupta
Decomposition

- Break up problem into tasks that can be carried out in parallel
  - Decomposition need not happen statically
  - New tasks can be identified as program executes

- Main idea: create at least enough tasks to keep all execution units on a machine busy

Key aspect of decomposition: identifying dependencies (or... a lack of dependencies)
Amdahl’s Law: dependencies limit maximum speedup due to parallelism

- You run your favorite sequential program...

- Let $S =$ the fraction of sequential execution that is inherently sequential (dependencies prevent parallel execution)

- Then maximum speedup due to parallel execution $\leq \frac{1}{S}$
A simple example

- Consider a two-step computation on a \( N \times N \) image
  - Step 1: double brightness of all pixels
    (independent computation on each grid element)
  - Step 2: compute average of all pixel values

- Sequential implementation of program
  - Both steps take \( \sim N^2 \) time, so total time is \( \sim 2N^2 \)
First attempt at parallelism (P processors)

- **Strategy:**
  - Step 1: execute in parallel
    - time for phase 1: $\frac{N^2}{P}$
  - Step 2: execute serially
    - time for phase 2: $N^2$

- **Overall performance:**
  
  Speedup $\leq \frac{2n^2}{n^2 / P + n^2}$
  
  Speedup $\leq 2$
Parallelizing step 2

- **Strategy:**
  - Step 1: execute in parallel
    - time for phase 1: \( \frac{N^2}{P} \)
  - Step 2: compute partial sums in parallel, combine results serially
    - time for phase 2: \( \frac{N^2}{P} + P \)

- **Overall performance:**
  - Speedup \( \leq \frac{2n^2}{\frac{2n^2}{p} + p} \)

Note: speedup \( \rightarrow P \) when \( N >> P \)
Amdahl’s law

- Let $S$ = the fraction of total work that is inherently sequential
- Max speedup on $P$ processors given by:

$$\text{speedup} \leq \frac{1}{1 - S} \frac{1}{S + \frac{1}{P}}$$

![Graph showing speedup vs. processors for different values of $S$ (0.01, 0.05, 0.1)]
Decomposition

- Who is responsible for performing decomposition?
  - In most cases: the programmer

- Automatic decomposition of sequential programs continues to be a challenging research problem (very difficult in general case)
  - Compiler must analyze program, identify dependencies
    - What if dependencies are data dependent (not known at compile time)?
  - Researchers have had modest success with simple loop nests
  - The “magic parallelizing compiler” for complex, general-purpose code has not yet been achieved
Assignment

** I had to pick a term
Assignment

- Assigning tasks to threads
  - Think of “tasks” as things to do
  - Think of threads as “workers”

- Goals: balance workload, reduce communication costs

- Can be performed statically, or dynamically during execution

- While programmer often responsible for decomposition, many languages/runtimes take responsibility for assignment.

** I had to pick a term (will explain in a second)
Assignment examples in ISPC

```c
export void sinx(
    uniform int N,
    uniform int terms,
    uniform float* x,
    uniform float* result)
{
    // assumes N % programCount = 0
    for (uniform int i=0; i<N; i+=programCount)
    {
        int idx = i + programIndex;
        float value = x[idx];
        float numer = x[idx] * x[idx] * x[idx];
        uniform int denom = 6;  // 3!
        uniform int sign = -1;
        for (uniform int j=1; j<=terms; j++)
        {
            value += sign * numer / denom;
            numer *= x[idx] * x[idx];
            denom *= (2*j+2) * (2*j+3);
            sign *= -1;
        }
        result[i] = value;
    }
}
```

Decomposition of work by loop iteration

**Programmer-managed assignment:**

*Static assignment*
Assign iterations to ISPC program instances in interleaved fashion

**System-managed assignment:**

*foreach construct exposes independent work to system*
System-manages assignment of iterations (work) to ISPC program instances (abstraction leaves room for dynamic assignment, but current ISPC implementation is static)

```c
export void sinx(
    uniform int N,
    uniform int terms,
    uniform float* x,
    uniform float* result)
{
    foreach (i = 0 ... N)
    {
        float value = x[i];
        float numer = x[i] * x[i] * x[i];
        uniform int denom = 6;  // 3!
        uniform int sign = -1;
        for (uniform int j=1; j<=terms; j++)
        {
            value += sign * numer / denom;
            numer *= x[i] * x[i];
            denom *= (2*j+2) * (2*j+3);
            sign *= -1;
        }
        result[i] = value;
    }
}
```
## Static assignment example using pthreads

```c
typedef struct {
    int N, terms;
    float* x, *result;
} my_args;

void parallel_sinx(int N, int terms, float* x, float* result)
{
    pthread_t thread_id;
    my_args args;

    args.N = N/2;
    args.terms = terms;
    args.x = x;
    args.result = result;

    // launch second thread, do work on first half of array
    pthread_create(&thread_id, NULL, my_thread_start, &args);

    // do work on second half of array in main thread
    sinx(N - args.N, terms, x + args.N, result + args.N);

    pthread_join(thread_id, NULL);
}

void my_thread_start(void* thread_arg)
{
    my_args* thread_args = (my_args*)thread_arg;
    sinx(thread_args->N, thread_args->terms, thread_args->x, thread_args->result); // do work
}
```

### Decomposition of work by loop iteration

**Programmer-managed assignment:**

**Static assignment**

Assign iterations to pthreads in blocked fashion (first half of array to spawned thread, second half to main thread)
Dynamic assignment using ISPC tasks

```c
void foo(uniform float* input,
         uniform float* output,
         uniform int N)
{
    // create a bunch of tasks
    launch[100] my_ispc_task(input, output, N);
}
```

**ISPC runtime assign tasks to worker threads**

**List of tasks:**

| task 0 | task 1 | task 2 | task 3 | task 4 | ... | task 99 |

**Assignment policy:** after completing current task, worker thread inspects list and assigns itself the next uncompleted task.
Orchestration

Problem to solve

Decomposition

Subproblems (a.k.a. “tasks”, “work to do”)

Assignment

Parallel Threads **
(“workers”)

Parallel program (communicating threads)

Orchestration

Mapping

Execution on parallel machine

** I had to pick a term
Orchestration

- **Involves:**
  - Structuring communication
  - Adding synchronization to preserve dependencies if necessary
  - Organizing data structures in memory
  - Scheduling tasks

- **Goals:** reduce costs of communication/sync, preserve locality of data reference, reduce overhead, etc.

- **Machine details impact many of these decisions**
  - If synchronization is expensive, might use it more sparsely
Mapping to hardware

Problem to solve

Decomposition

Subproblems
(a.k.a. “tasks”, “work to do”)

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(“workers”)

Parallel program
(communicating threads)

Execution on parallel machine

Assignment

Orchestration

Mapping

** I had to pick a term
Mapping to hardware

- Mapping “threads” (“workers”) to hardware execution units

- **Example 1: mapping by the operating system**
  - e.g., map pthread to HW execution context on a CPU core

- **Example 2: mapping by the compiler**
  - Map ISPC program instances to vector instruction lanes

- **Example 3: mapping by the hardware**
  - Map CUDA thread blocks to GPU cores (future lecture)

- **Some interesting mapping decisions:**
  - Place related threads (cooperating threads) on the same processor (maximize locality, data sharing, minimize costs of comm/sync)
  - Place unrelated threads on the same processor (one might be bandwidth limited and another might be compute limited) to use machine more efficiently
Decomposing computation or data?

Often, the reason a problem requires lots of computation (and needs to be parallelized) is that it involves manipulating a lot of data.

I’ve described the process of parallelizing programs as an act of partitioning **computation** (work).

Often, it’s equally valid to think of **partitioning data**. (computations go with the data)

But there are many computations where the correspondence between work-to-do (”tasks”) and data is less clear. In these cases it’s natural to think of partitioning computation.
A parallel programming example
A 2D-grid based solver

- Solve partial differential equation (PDE) on \((N+2) \times (N+2)\) grid
- Iterative solution
  - Perform Gauss-Seidel sweeps over grid until convergence

\[
\]
Grid solver algorithm
C-like pseudocode for sequential algorithm is provided below

```c
const int n;
float* A; // assume allocated to grid of N+2 x N+2 elements
void solve(float* A) {
  float diff, prev;
  bool done = false;
  while (!done) { // outermost loop: iterations
    diff = 0.f;
    for (int i=1; i<n; i++) { // iterate over non-border points of grid
      for (int j=1; j<n; j++) {
        prev = A[i,j];
                         A[i,j+1] + A[i+1,j]);
        diff += fabs(A[i,j] - prev); // compute amount of change
      }
    }
    if (diff/(n*n) < TOLERANCE) // quit if converged
      done = true;
  }
}
```

Grid solver example from: Culler, Singh, and Gupta
Step 1: identify dependencies (problem decomposition phase)

Each row element depends on element to left.
Each column depends on previous column.
Step 1: identify dependencies (problem decomposition phase)

There is independent work along the diagonals!

Good: parallelism exists!

Possible implementation strategy:
1. Partition grid cells on a diagonal into tasks
2. Update values in parallel
3. When complete, move to next diagonal

Bad: independent work is hard to exploit
Not much parallelism at beginning and end of computation.
Frequent synchronization (after completing each diagonal)
Let’s make life easier on ourselves

- Idea: improve performance by changing the algorithm to one that is more amenable to parallelism
  - Change the order grid cell cells are updated
  - New algorithm iterates to same solution (approximately), but converges to solution differently
    - Note: floating-point values computed are different, but solution still converges to within error threshold
  - Yes, we needed domain knowledge of Gauss-Seidel method for solving a linear system to realize this change is permissible for the application
New approach: reorder grid cell update via red-black coloring

Update all red cells in parallel

When done updating red cells, update all black cells in parallel (respect dependency on red cells)

Repeat until convergence
Possible assignments of work to processors

Question: Which is better? Does it matter?
Answer: it depends on the system this program is running on
Consider dependencies (data flow)

1. Perform red update in parallel
2. Wait until all processors done with update
3. Communicate updated red cells to other processors
4. Perform black update in parallel
5. Wait until all processors done with update
6. Communicate updated black cells to other processors
7. Repeat
Communication resulting from assignment

** Blocked Assignment **

** Interleaved Assignment **

= data that must be sent to P2 each iteration

Blocked assignment requires less data to be communicated between processors
Data-parallel expression of solver
Data-parallel expression of grid solver

Note: to simplify pseudocode: just showing red-cell update

```c
const int n;

float* A = allocate(n+2, n+2));  // allocate grid

void solve(float* A) {
  bool done = false;
  float diff = 0.f;
  while (!done) {
    for_all (red cells (i,j)) {
      float prev = A[i,j];
      reduceAdd(diff, abs(A[i,j] - prev));
    }
    if (diff/(n*n) < TOLERANCE)
      done = true;
  }
}
```

Data-parallel expression of grid solver

Note: to simplify pseudocode: just showing red-cell update

Assignment: ???

Orchestration: handled by system
(builtin communication primitive: reduceAdd)

Orchestration: handled by system
(End of for_all block is implicit wait for all workers before returning to sequential control)

Grid solver example from: Culler, Singh, and Gupta
Shared address space (with SPMD threads)
expression of solver
Shared address space expression of solver

SPMD execution model

- Programmer is responsible for synchronization
- Common synchronization primitives:
  - Locks (provide mutual exclusion): only one thread in the critical region at a time
  - Barriers: wait for threads to reach this point
Shared address space solver (pseudocode in SPMD execution model)

```c
int n;  // grid size
bool done = false;
float diff = 0.0;
LOCK myLock;
BARRIER myBarrier;

// allocate grid
float* A = allocate(n+2, n+2);

void solve(float* A) {
    int threadId = getThreadId();
    int myMin = 1 + (threadId * n / NUM_PROCESSORS);
    int myMax = myMin + (n / NUM_PROCESSORS);

    while (!done) {
        diff = 0.f;
        barrier(myBarrier, NUM_PROCESSORS);
        for (j=myMin to myMax) {
            for (i = red cells in this row) {
                float prev = A[i,j];
                lock(myLock)
                diff += abs(A[i,j] - prev);
                unlock(myLock);
            }
        }
        barrier(myBarrier, NUM_PROCESSORS);
        if (diff/(n*n) < TOLERANCE)    // check convergence, all threads get same answer
            done = true;
        barrier(myBarrier, NUM_PROCESSORS);
    }
}
```

Assume these are global variables (accessible to all threads)

Assume solve function is executed by all threads. (SPMD-style)

Value of threadId is different for each SPMD instance: use value to compute region of grid to work on

Each thread computes the rows it is responsible for updating

Grid solver example from: Culler, Singh, and Gupta
Review: need for mutual exclusion

- Each thread executes
  - Load the value of diff into register r1
  - Add the register r2 to register r1
  - Store the value of register r1 into diff

- One possible interleaving: (let starting value of diff=0, r2=1)

<table>
<thead>
<tr>
<th>T0</th>
<th>T1</th>
</tr>
</thead>
<tbody>
<tr>
<td>r1 ← diff</td>
<td>T0 reads value 0</td>
</tr>
<tr>
<td>r1 ← r1 + r2</td>
<td>T1 reads value 0</td>
</tr>
<tr>
<td>diff ← r1</td>
<td>T0 sets value of its r1 to 1</td>
</tr>
<tr>
<td></td>
<td>T1 sets value of its r1 to 1</td>
</tr>
<tr>
<td></td>
<td>T0 stores 1 to diff</td>
</tr>
<tr>
<td></td>
<td>T1 stores 1 to diff</td>
</tr>
</tbody>
</table>

- Need this set of three instructions to be atomic
Mechanisms for preserving atomicity

- Lock/unlock mutex around a critical section
  
  LOCK(mylock);
  // critical section
  UNLOCK(mylock);

- Some languages have first-class support for atomicity of code blocks
  
  atomic {
    // critical section
  }

- Intrinsics for hardware-supported atomic read-modify-write operations
  
  atomicAdd(x, 10);
Shared address space solver (pseudocode in SPMD execution model)

```c
int n;       // grid size
bool done = false;
float diff = 0.0;
LOCK myLock;
BARRIER myBarrier;

// allocate grid
float* A = allocate(n+2, n+2);

void solve(float* A) {
    int threadId = getThreadId();
    int myMin = 1 + (threadId * n / NUM_PROCESSORS);
    int myMax = myMin + (n / NUM_PROCESSORS);

    while (!done) {
        diff = 0.f;
        barrier(myBarrier, NUM_PROCESSORS);
        for (j=myMin to myMax) {
            for (i = red cells in this row) {
                float prev = A[i,j];
                                A[i+1,j], A[i,j+1]);
                lock(myLock)
                diff += abs(A[i,j] - prev));
                unlock(myLock);
            }
        }
        barrier(myBarrier, NUM_PROCESSORS);
        if (diff/(n*n) < TOLERANCE) // check convergence, all threads get same answer
            done = true;
        barrier(myBarrier, NUM_PROCESSORS);
    }
}
```

Do you see a potential performance problem with this implementation?

Grid solver example from: Culler, Singh, and Gupta
Shared address space solver (SPMD execution model)

int n; // grid size
bool done = false;
float diff = 0.0;
LOCK myLock;
BARRIER myBarrier;

// allocate grid
float* A = allocate(n+2, n+2);

void solve(float* A) {
  float myDiff;
  int threadId = getThreadId();
  int myMin = 1 + (threadId * n / NUM_PROCESSORS);
  int myMax = myMin + (n / NUM_PROCESSORS);

  while (!done) {
    float myDiff = 0.f;
    diff = 0.f;
    barrier(myBarrier, NUM_PROCESSORS);
    for (j=myMin to myMax) {
      for (i = red cells in this row) {
        float prev = A[i,j];
          A[i+1,j], A[i,j+1]);
        myDiff += abs(A[i,j] - prev));
      }
      lock(myLock);
      diff += myDiff;
      unlock(myLock);
      barrier(myBarrier, NUM_PROCESSORS);
      if (diff/(n*n) < TOLERANCE) // check convergence, all threads get same answer
        done = true;
    }
  }
}

Improve performance by accumulating into partial sum locally, then complete reduction globally at the end of the iteration.

Now only only lock once per thread, not once per (i,j) loop iteration!

Grid solver example from: Culler, Singh, and Gupta
Barrier synchronization primitive

- `barrier(num_threads)`
- Barriers are a conservative way to express dependencies
- Barriers divide computation into phases
- All computations by all threads before the barrier complete before any computation in any thread after the barrier begins
int n; // grid size
bool done = false;
float diff = 0.0;
LOCK myLock;
BARRIER myBarrier;

// allocate grid
float* A = allocate(n+2, n+2);

void solve(float* A) {
    float myDiff;
    int threadId = getThreadId();
    int myMin = 1 + (threadId * n / NUM_PROCESSORS);
    int myMax = myMin + (n / NUM_PROCESSORS);

    while (!done) {
        float myDiff = 0.f;
        diff = 0.f;
        barrier(myBarrier, NUM_PROCESSORS);
        for (j=myMin to myMax) {
            for (i = red cells in this row) {
                float prev = A[i,j];
                                A[i+1,j], A[i,j+1]);
                myDiff += abs(A[i,j] - prev));
            }
        } 
        lock(myLock);
        diff += myDiff;
        unlock(myLock);
        barrier(myBarrier, NUM_PROCESSORS); // check convergence, all threads get same answer
        if (diff/(n*n) < TOLERANCE)
            done = true;
        barrier(myBarrier, NUM_PROCESSORS);
    }
}

Grid solver example from: Culler, Singh, and Gupta
Grid solver example from: Culler, Singh, and Gupta

```c
int n;  // grid size
bool done = false;
LOCK myLock;
BARRIER myBarrier;
float diff[3]; // global diff, but now 3 copies

float *A = allocate(n+2, n+2);

void solve(float* A) {
    float myDiff;  // thread local variable
    int index = 0;  // thread local variable

    diff[0] = 0.0f;
    barrier(myBarrier, NUM_PROCESSORS);  // one-time only: just for init

    while (!done) {
        myDiff = 0.0f;
        // // perform computation (accumulate locally into myDiff)
        // lock(myLock);
        diff[index] += myDiff;  // atomically update global diff
        unlock(myLock);
        diff[(index+1) % 3] = 0.0f;
        barrier(myBarrier, NUM_PROCESSORS);
        if (diff[index]/(n*n) < TOLERANCE)
            break;
        index = (index + 1) % 3;
    }
}
```

Idea:

Remove dependencies by using different diff variables in successive loop iterations

Trade off footprint for removing dependencies!
(a common parallel programming technique)
More on specifying dependencies

- Barriers: simple, but conservative (coarse-granularity dependencies)
  - All work in program up until this point (for all threads) must finish before any thread begins next phase

- Specifying specific dependencies can increase performance (by revealing more parallelism)
  - Example: two threads. One produces a result, the other consumes it.

```
// produce x, then let T1 know
x = 1;
flag = 1;
// do more work here...

// do stuff independent
// of x here
while (flag == 0);
print x;
```

- We just implemented a message queue (of length 1)
Solver implementation in two programming models

- Data-parallel programming model
  - Synchronization:
    - Single logical thread of control, but iterations of `for all` loop may be parallelized by the system (implicit barrier at end of `for all` loop body)
  - Communication
    - Implicit in loads and stores (like shared address space)
    - Special built-in primitives for more complex communication patterns: e.g., reduce

- Shared address space
  - Synchronization:
    - Mutual exclusion required for shared variables (e.g., via locks)
    - Barriers used to express dependencies (between phases of computation)
  - Communication
    - Implicit in loads/stores to shared variables
We will defer discussion of the message passing expression of solver to a later class.
Summary

- **Amdahl’s Law**
  - Overall maximum speedup from parallelism is limited by amount of serial execution in a program

- **Aspects of creating a parallel program**
  - Decomposition to create independent work, assignment of work to workers, orchestration (to coordinate processing of work by workers), mapping to hardware
  - We’ll talk a lot about making good decisions in each of these phases in the coming lectures (in practice, they are very inter-related)

- **Focus today: identifying dependencies**

- **Focus soon: identifying locality, reducing synchronization**